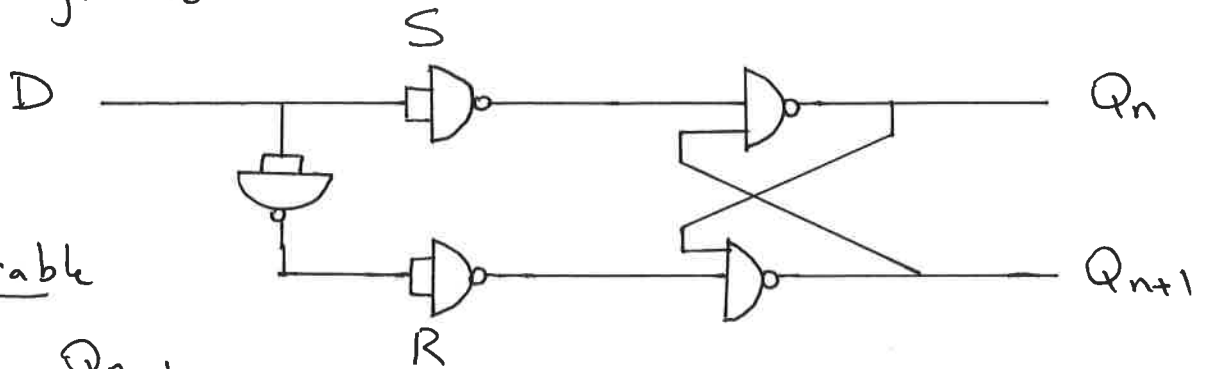


The Logic cet:-



ch/s table

| Q_n | D | Q_{n+1} |
|-------|---|-----------|
| 0 | 0 | 0 |
| 0 | 1 | 1 |
| 1 | 0 | 0 |
| 1 | 1 | 1 |

| $Q_n \backslash D$ | 0 | 1 |
|--------------------|---|---|
| 0 | 0 | 1 |
| 1 | 2 | 1 |

$Q_{n+1} = D$ ch/s eq

T Flip-Flop :-

The trigger F.F has one input, T, such that if $T=1$, the F.F changes state (that is, trigger), and if $T=0$, the state remains the same.

The truth table of the T/F.F can shown below:-

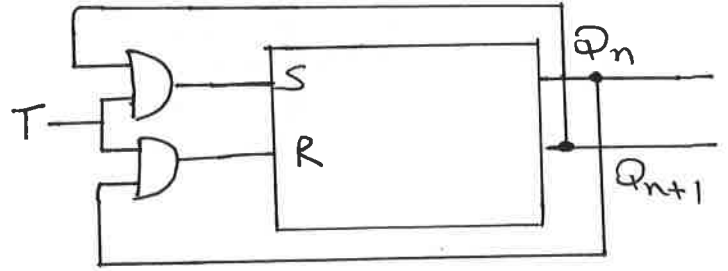
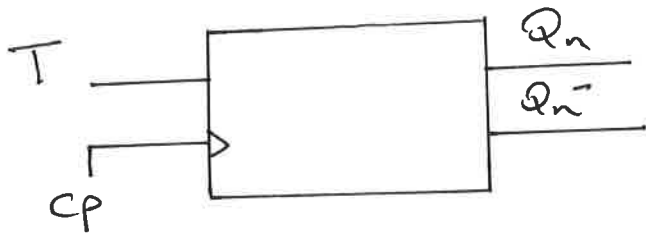
| I/P T | present Q_n | T/F.F next Q_{n+1} |
|----------|------------------|----------------------------|
| 0 | 0 | 0 |
| 0 | 1 | 1 |
| 1 | 0 | 1 |
| 1 | 1 | 0 |

* This means $(Q_{n+1}) = Q_n$ when $T=0$

* $(Q_{n+1}) = \bar{Q}_n$, when $T=1$

$Q_{n+1} = T \oplus Q_n$

The Block Diagram and the cct Diagram of T/F.F is:-



if

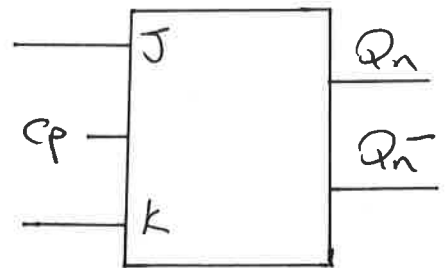
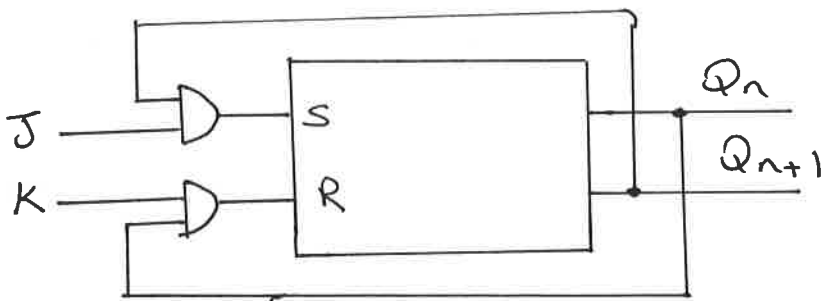
| T | Q_n | Q_{n+1} |
|---|-------|-----------|
| 1 | 0 | 1 |

set(1)

| T | Q_n | Q_{n+1} |
|---|-------|-----------|
| 1 | 1 | 0 |

J-K Flip-Flop:-

It is the most commonly used type of F.F. It is very similar to SR/F.F except that J.K/F.F allow to used Logic 1 on both inputs J & K, the cct Diagram of J.K/F.F is shown below:-



And the truth table of this type is shown below:-

| input | | Present state (Q_n) | Next state (Q_{n+1}) |
|-------|---|-------------------------|--------------------------|
| J | K | | |
| 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 1 |
| 0 | 1 | 0 | 0 |
| 0 | 1 | 1 | 0 |
| 1 | 0 | 0 | 1 |
| 1 | 0 | 1 | 1 |
| 1 | 1 | 0 | 1 |
| 1 | 1 | 1 | 0 |

This table can be put in another form such that:-

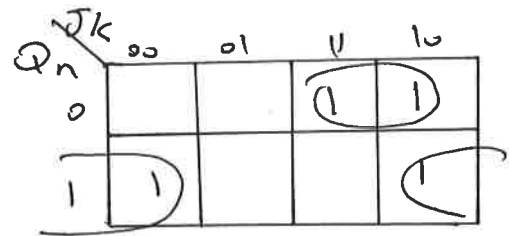
| J | K | Q_{n+1} |
|---|---|-------------|
| 0 | 0 | Q_n |
| 0 | 1 | 0 |
| 1 | 0 | 1 |
| 1 | 1 | \bar{Q}_n |

From the set Diagram, we can conclude that:-

$$S = J \bar{Q}_n, \quad R = K Q_n$$

The ch/s equation of JK/F.F can be:-

$$Q_{n+1} = J \bar{Q}_n + \bar{K} Q_n$$

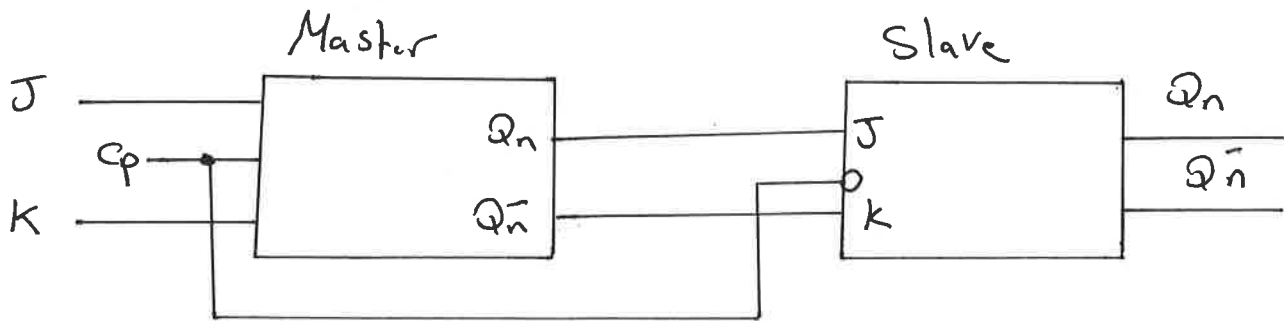


J-K master-slave F/F :-

A master-slave F.F is a combination of 2 clocked F.F J.K or S-R/F.F.

The first is called "The master" and the second "The slave" notes that the master is positively clocked but the slave is negative.

The Block Diagram of such Flip-Flop is shown below



Flip-Flop Excitation table :-

The tables of the four standard Flip-Flops can be

summarized as :-

* S-R / F.F

| S | R | Q_{n+1} |
|---|---|-----------|
| 0 | 0 | Q_n |
| 0 | 1 | 0 |
| 1 | 0 | 1 |
| 1 | 1 | X |

* J-K / F.F

| J | K | Q_{n+1} |
|---|---|-----------|
| 0 | 0 | Q_n |
| 0 | 1 | 0 |
| 1 | 0 | 1 |
| 1 | 1 | Q_n |

* D / F.F

| D | Q_{n+1} |
|---|-----------|
| 0 | 0 |
| 1 | 1 |

* T/F.F

| T | Q_n | Q_{n+1} |
|---|-------|-----------|
| 0 | 0 | 0 |
| 0 | 1 | 1 |
| 1 | 0 | 1 |
| 1 | 1 | 0 |

The Excitation table of the same F/F gives the relation between Q_n and Q_{n+1} as input and the standard i/p as o/p.

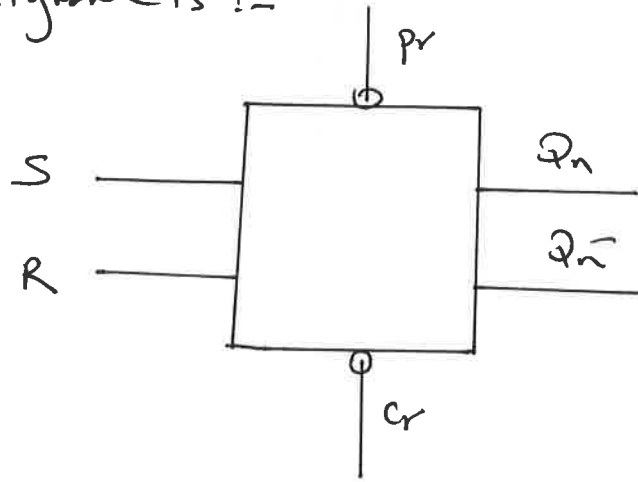
| Q_n | Q_{n+1} | S | R | J | K | D | T |
|-------|-----------|---|---|---|---|---|---|
| 0 | 0 | 0 | d | 0 | d | 0 | 0 |
| 0 | 1 | 1 | 0 | 1 | d | 1 | 1 |
| 1 | 0 | 0 | 1 | d | 1 | 0 | 1 |
| 1 | 1 | d | 0 | d | 0 | 1 | 0 |

A synchronous inputs:-

The S-R inputs or (J-K) are called synchronous i/p's since they only effect the state of the F.F when the clock pulse is present. Some F.F have either i/p's called preset or set and clear or reset that can be used for setting F.F to Logic 1 or resetting to Logic zero, by applying the proper signal to the preset and clear inputs.

These inputs can change the state of the F.F regardless of the synchronous i/p of the block.

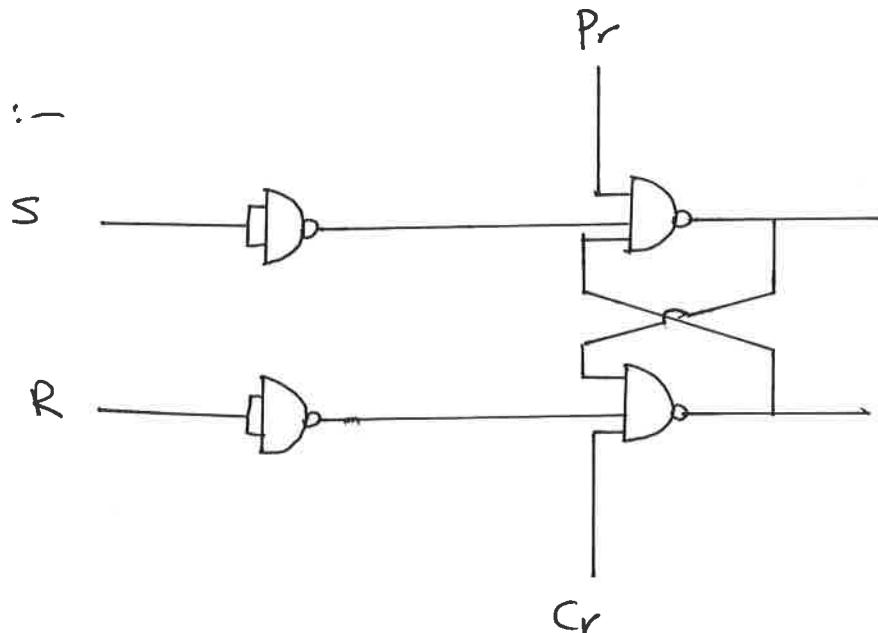
The Block Diagram is :-



The (Pr and Cr) can be used to determine the primary state of the Flip-Flops or/for we can setting the F.F at a certain value, regardless of input values of (S, R, J, K), and can put the table as:-

| Cr | Pr | Condition (Qn) |
|----|----|------------------|
| 0 | 0 | Don't use |
| 0 | 1 | 0 |
| 1 | 0 | 1 |
| 1 | 1 | Normal condition |

The Logic cct:-

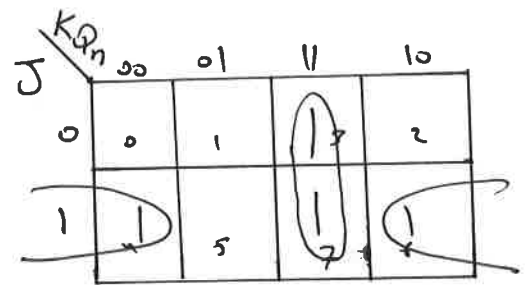


Transformation of Flip-Flops:-

* Transform T/F.F to J-K/F.F

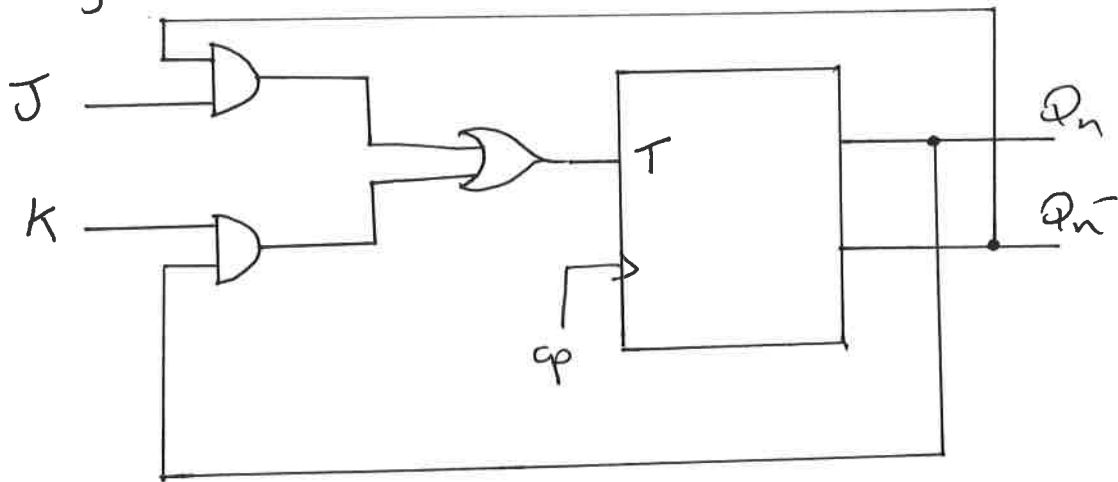
* Truth table of J-K/F.F

| J | K | Q_n | Q_{n+1} | T |
|---|---|-------|-----------|---|
| 0 | 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 1 | 0 |
| 0 | 1 | 0 | 0 | 0 |
| 0 | 1 | 1 | 0 | 1 |
| 1 | 0 | 0 | 1 | 1 |
| 1 | 0 | 1 | 1 | 0 |
| 1 | 1 | 0 | 1 | 1 |
| 1 | 1 | 1 | 0 | 1 |



* ch/s eq. $T = KQ_n + JQ_n$

* The Logic cct:-



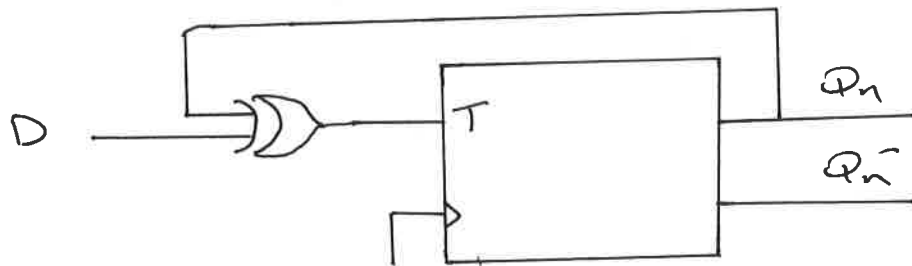
Ex) Transform T/F.F to D/F.F

| D | Q_n | Q_{n+1} | T |
|---|-------|-----------|---|
| 0 | 0 | 0 | 0 |
| 0 | 1 | 0 | 1 |
| 1 | 0 | 1 | 1 |
| 1 | 1 | 1 | 0 |

| Q_n \ D | 0 | 1 |
|-----------|---|---|
| 0 | | 1 |
| 1 | 1 | |

chls eq $T = \bar{D}Q_n + DQ_n^-$
 $= D \oplus Q_n$

Logic ckt:-



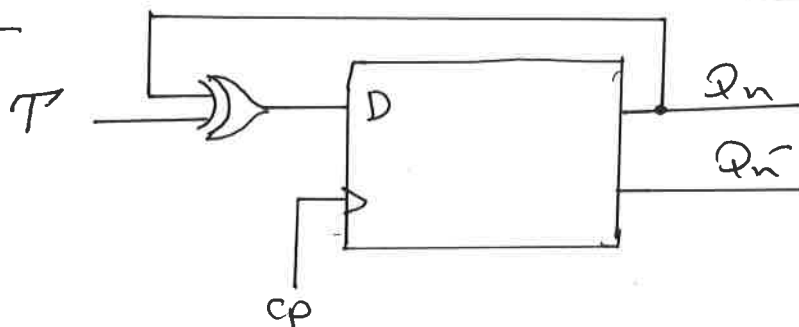
Ex) Transform D/F.F to T/F.F

| T | Q_n | Q_{n+1} | D |
|---|-------|-----------|---|
| 0 | 0 | 0 | 0 |
| 0 | 1 | 1 | 1 |
| 1 | 0 | 1 | 1 |
| 1 | 1 | 0 | 0 |

| Q_n \ T | 0 | 1 |
|-----------|---|---|
| 0 | | 1 |
| 1 | 1 | |

$D = \bar{T}Q_n + TQ_n^-$
 $= T \oplus Q_n$

Logic ckt:-



Ex) convert S-R/F.F to D/F.F

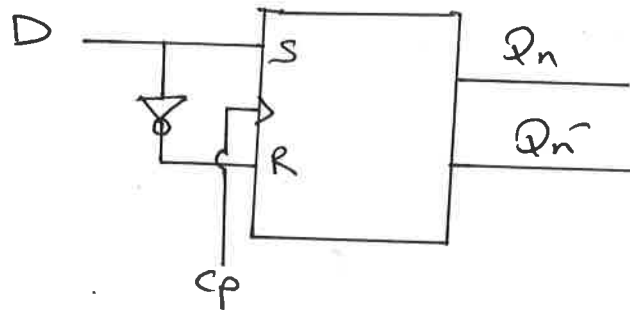
| D | Q_n | Q_{n+1} | S | R |
|---|-------|-----------|---|---|
| 0 | 0 | 0 | 0 | d |
| 0 | 1 | 0 | 0 | 1 |
| 1 | 0 | 1 | 1 | 0 |
| 1 | 1 | 1 | d | 0 |

| Q_n | D | 0 | 1 |
|-------|---|---|---|
| 0 | 0 | 0 | 1 |
| 1 | 0 | 0 | d |

$S = D$

| Q_n | D | 0 | 1 |
|-------|---|---|---|
| 0 | d | 0 | 0 |
| 1 | 1 | 0 | 0 |

$R = \bar{D}$



Ex) Design T/F.F by using S-R/F.F?

| T | Q_n | Q_{n+1} | S | R |
|---|-------|-----------|---|---|
| 0 | 0 | 0 | 0 | d |
| 0 | 1 | 1 | d | 0 |
| 1 | 0 | 1 | 1 | 0 |
| 1 | 1 | 0 | 0 | 1 |

| Q_n | T | 0 | 1 |
|-------|---|---|---|
| 0 | 0 | 0 | 1 |
| 1 | d | 0 | 0 |

$S = TQ_n$

| Q_n | T | 0 | 1 |
|-------|---|---|---|
| 0 | d | 0 | 1 |
| 1 | 0 | 0 | 1 |

$R = TQ_n$

